Semantic Web 0 (0) 1 IOS Press

OPTIMA: A Hybrid OBDA System for Efficiently Querying Large Heterogeneous Data

Chahrazed B. Bachir Belmehdi^{a,*}, Abderrahmane Khiat^b and Nabil Keskes^a

^a Department first, ESI-SBA Institute, SBA, Algeria

E-mail: cb.bachirbelmehdi@esi-sba.dz

^b Enterprise Information Systems, Fraunhofer IAIS, Bonn, Germany

E-mails: abderrahmane.khiat@iais.fraunhofer.de, n.keskes@esi-sba.dz

Abstract. The current decade is witnessing a remarkable evolution in terms of big data virtualization. Data is queried on-the-fly against the original data sources without any prior data materialization. Ontology-Based Big Data Access solutions by design use a fixed model, e.g., TABULAR, as the only Virtual Data Model - a uniform schema that is built on-the-fly to load, transform, and join relevant data. While other data models such as GRAPH or DOCUMENT are more flexible and, thus, can be more suitable for some common types of queries such as join or nested queries. Those queries are, in many cases, hard to predict because they depend on many criteria such as query plan, data model, data size, operations e.g., join, filter. To address the problem of selecting the optimal virtual data model for various queries on large datasets, we develop OPTIMA. OPTIMA is a framework that (1) builds on the principle of ontology-based data access to enable the querying, aggregating, and joining of large heterogeneous data in a distributed manner using a unique query language SPARQL and (2) calls the deep learning method to predict the optimal virtual data model using the features extracted from SPARQL queries. OPTIMA currently leverages state-of-the-art Big Data technologies, Spark, and implements two virtual data models, GRAPH and TABULAR, and supports out-of-the-box five data sources Neo4j, MongoDB, MySQL, Cassandra, and CSV. Extensive experiments show that OPTIMA returns the optimal virtual model with an accuracy of 0.831, thus reducing the query execution time by over 40% in favor of tabular model selection and over 30% for the graph model selection.

OPTIMA is available on GitHub https://github.com/chahrazedbb/OPTIMA

Keywords: Data Virtualization, Big Data, Ontology Based Data Access, Deep Learning

1. Introduction

Data virtualization approaches tackle data integration challenges by creating a virtual data model under which the heterogeneous formats are homogenized *on-the-fly* without data materialization [1]. Ontology-Based Data Access (OBDA) [2] approaches [3] maintain data virtualization with practical knowledge representation models and ontology-based mappings. However, existing solutions dedicated to big data [4][5][6] use by design a fixed model, e.g., TABULAR as the only virtual data model¹ to load and transform the requested data into a uniform model to be joined.

*Corresponding author. E-mail: cb.bachirbelmehdi@esi-sba.dz.

¹We denote GRAPH and TABULAR to distinguish between Virtual model and data source models.

N. Bachir Belmehdi et al. / OPTIMA



Fig. 1. OPTIMA Framework

Nevertheless, the TABULAR virtual model can have performance downsides for queries involving many join operations on large data [7]. In contrast, other data models, such as GRAPH, perform better for such queries. On the other hand, the TABULAR model performs better for queries that involve selection, or projection [8]. Therefore, there is a need to support different virtual models and select the optimal one based on query behavior, thus saving operational execution time. However, predicting the optimal virtual model is challenging since the selection depends on many criteria, such as query plan, data model, size, and operations.

This challenge raises the following research questions: based on a given query, RQ1: does the data model affect query execution time? RQ2: which virtual data model is optimal, i.e., the model that has the lowest cost, precisely, the lowest query execution time? and how to select it?"

To address this research problem, we introduce our operational tool OPTIMA, which was briefly presented and published in our earlier work [9]. OPTIMA is an extensible framework that uses two Virtual models, GRAPH and TABULAR, and supports out-of-the-box five data sources, Property Graph, Relational, Tabular, Document-based, and Wide-Columnar, to query large heterogeneous data in an efficient way. To select the optimal virtual data model GRAPH or TABULAR, we used one hot vector encoding to transform different SPARQL features into hidden representations. Next, we embed these representations into a tree-structured model, which we use to classify the virtual model GRAPH or TABULAR that has the lowest query execution time. Experiments show that OPTIMA reduces the query execution time of over 40% for the TABULAR model selection and over 30% for the GRAPH model selection. We describe each component of our system OPTIMA as illustrated in Figure 1, Label 1; followed by conducted experiment and related work.

2. OPTIMA: Optimal Virtual Model for Querying Large Heterogeneous Data

2.1. Virtual Data Model Prediction

Built on top of OBDA components, this distinctive component implemented in OPTIMA aims to select the optimal virtual data model GRAPH or TABULAR based on the query behavior see Figure 2. The component receives the SPARQL query as input and predicts the optimal virtual data model with the lowest execution time. The deep learning model starts by breaking down the SPARQL query plan into nodes. Each node includes a set of query features that significantly affect the query execution time (e.g., filter). The different features are then encoded using a one-hot vector. Next, we propose a tree-structured model that takes the encoded features of SPARQL query as input to learn each sub-plan representation effectively and outputs the optimal virtual data model, GRAPH or TABULAR,

that has the lowest query execution time. Our model consists of an embedding layer to condense the features' vectors and an estimation layer to estimate the optimal virtual data model. In addition, the model includes an intermediate representation layer to capture the correlation between the joined star-shaped queries. Once the optimal model is predicted, the rest of the OBDA components and operations (e.g., join) follow the optimal virtual model predicted

GRAPH or TABULAR.



Fig. 2. Architecture of Predictive Model for Optimal Virtual Data Model

2.2. Query Decomposition and Relevant Entity Detection

This component decomposes the SPARQL query into star-shaped queries. More precisely, the query's Basic Graph Pattern (BGP) is divided into a set of sub-BGPs, where each sub-BGP contains all the triple patterns sharing the same subject variable. Those sub-BGPs sharing the same subject are called a star-shaped query. Next, this component analyzes each star-shaped query and visits the mappings file to obtain the data source's path and the attributes mapped to each element of the star-shaped query, i.e., relevant entities. This information is then passed to the data wrapper to load relevant entities.

2.3. Data Wrapper

Once the sources and relevant entities are identified using mappings, the data wrapper converts relevant entities (e.g., tables) from their original models to data that comply with the optimal virtual data model predicted, which is actually the data structure of the computation unit of the query engine. This conversion occurs at query-time, which allows for the parallel execution of expensive operations, e.g., join. Query engines implement wrappers called connectors to convert data entities from the source to the virtual data model, performing transformation of data source, e.g., relational model to virtual model, e.g., GRAPH (see Figure 3a).

Each star-shaped query corresponds to one relevant entity, and thus one single virtual data model is created. According to the mapping, this occurs when the relevant entity is retrieved only from one data source. Otherwise, if the relevant entity is retrieved from multiple data sources, then the virtual model for one relevant entity is the union of the temporary virtual models created for each source (see Figure 3b). Below we describe the data sources' model transformation by wrappers into GRAPH and TABULAR.

- For the virtual data model of type GRAPH, one relevant entity of a star-shaped query from relational and tabular models is a table with specific columns, which is transformed into one virtual model GRAPH. Figure 3a, Label 1 illustrates the transformation process from relational into GRAPH, following the process described in [10]. For each table row, a vertex is created with the same label as the table's name (e.g., table' Person' corresponds to all vertices with the label "Person") in addition to the root vertex. Edges are created between vertices and the root vertex, whereas the properties of each vertex are the columns of the table (e.g., column 'name' corresponds to property 'name'), and the values of the properties are the table's cell information. The process is applied to a property graph that has the same data structure as the GRAPH. Thus, the transformation process is a direct mapping; the node corresponds to a vertex, the node's properties correspond to the properties of the vertex, and



wide-columnar is parsed to create a virtual TABULAR. As described in [12], the Virtual TABULAR consists of a table with a name similar to the root object's name (e.g., a table 'Person' from object name 'Person'). A new row is inserted by iterating through the object's elements into the corresponding table. The corresponding key-values are saved under the column representing the cell information (see Figure 3a, Label 2). As for the property graph, a Virtual TABULAR is created for each distinct graph that matches the pattern queried, following the approach proposed in [11]. The Virtual TABULAR consists of a table with the same name as the label shared by vertices. A default column' ID' is created to store each vertex id with the same label. A new row is inserted for each vertex into the corresponding table; For each distinct property of the vertex, an additional column is created with the same datatype as the extracted property. The cell information consists of the values extracted from the vertex's properties. Finally, the relational and tabular data structure is the same to some extent as the data structure of the Virtual TABULAR Model. Therefore, the transformation process is a direct mapping between both models.

 N. Bachir Belmehdi et al. / OPTIMA



2.4. Distributed Query Processor

Distributed Query Processor is the environment where queries are executed. OPTIMA calls for Graphx and Apache Spark to use two different virtual data models, GRAPH and TABULAR. We consider two types of data models, GRAPH and TABULAR, which allow for (1) graph-parallel computation (Figure 6a) and (2) data-parallel computation (Figure 6b), which affect the query performance. We should point out that we did not employ any query optimization function to choose the most efficient query execution plan; instead, we focused on the join operation. If deep learning predicts that the optimal virtual model is of type GRAPH, then for each relevant entity, one virtual GRAPH model is generated by wrappers. The wrappers use API to access the data source and perform the transformation. OPTIMA joins those GRAPHs into a Final Virtual GRAPH (see Figure Figure 4a) using "multi-join algorithm" (see Figure Figure 4b, Label 1) or TABULARs into a Final Virtual TABULAR using "incremental join algorithm" (see Figure Figure 4b, Label 2). This joining is through connections between star-shaped queries; see Figure 1 Label 2 and Figure 5. However, GRAPH and TABULAR have different structures. For example, the interaction with GRAPH is possible by means of Graph Pattern Matching operations (Cypher-like), while the interaction with TABULAR is possible by SQL-like functions. SPARQL and star-shaped query operations (e.g., limit) are translated into Virtual Data model operations (e.g., "take" in the case of Graphx).

3. Experimental Setup

We conducted an experimental study to evaluate OPTIMA performance compared to the state-of-the-art SPARKbased Sequerall, which uses dataframes (i.e., TABULAR) as a virtual data model. We used five tables to enable up to 44 4-chain joins. These tables are loaded in five different data sources Cassandra, MongoDB, CSV, Neo4j, and MySQL. 45 Table 1 shows the described information about data. We generated 5150 queries with 0-4 joins,0-45 selection, 0-16 46 filter, limit, and OrderBy. We take 4120 queries for training the model and 1030 queries for validation. We run the 47 evaluation on Ubuntu 64-bit with an Intel(R) Core(TM) i7-8550U CPU @ 1.80GHz, allocating 8GB of RAM.

We conducted an empirical study to evaluate OPTIMA performance and address our research problem with the following questions: RQ1: What is the query performance using OPTIMA? RQ2: Is the time of prediction plus the time of query execution using an optimal virtual model equal to the fixed one? RQ3: What is the query performance when using TABULAR versus GRAPH? RQ4: What is the accuracy of OPTIMA and machine learning? RQ5:

	Q1	Q2	Q3	Q4	Q5	Q6	Q7	Q8	Q9	Q10	Q11	Q12	Q13	Q14	Q15	Q16	Q17	Q18	Q19	Q20
Product	1	 Image: A set of the set of the		1	1	 Image: A set of the set of the	 Image: A set of the set of the	1	 Image: A set of the set of the	 Image: A set of the set of the			 Image: A set of the set of the	1	1	1	√	 Image: A set of the set of the	√	 Image: A second s
Offer	1		1	1	√	~	1	1	√	1	1			1	1	1	×		√	 Image: A second s
Review	1	 Image: A second s	1	1		 Image: A second s	1	1		1			1	1	1			1	1	 Image: A second s
Person			1	1		 Image: A second s			 Image: A second s		1	1						1	1	1
Producer	1		1	1		 Image: A second s	1	1			1				1			1	1	1
PROJECT	√ 16	√5	✓29	√ 45	✓24	√ 45	√38	√38	✓24	√34	√ 4	√ 6	✓32	√34	√ 4	√5	√ 9	√45	√ 45	√5
FILTER	√ 16		√ 12	√ 1		√5	√ 1	√ 1					√ 1	√ 1	√ 4			✓2	√ 3	
ORDERBY	√ 1	✓1			✓1		✓1		✓1	✓1			✓1	√ 1			✓1			
LIMIT	√300	✓2		✓20	✓4	✓20	✓20	√ 80			√ 10		√ 13	√ 19	√ 1000	√ 1000				

Tables and	Operations	involved	in	Oueries
rables and	operations	mvorvcu	111	Queries.

What is the query performance of OPTIMA compared to the state-of-the-art, e.g., Squerall [6]? RQ6: What is the impact of involving more data sources in a join query? RQ7: What is the resource consumption (CPU, memory) of OPTIMA while running various queries? RQ8: What is the time taken by each transformation process?

3.1. Method

We consider two studies:

- In the first study, we compare OPTIMA's results with SPARK-based Squerall's results. Our comprehensive literature review did not reveal any single work except Squerall that is available and that support most data sources. Squerall uses two big data engines, Presto and SPARK: Presto-based, where the virtual model of Presto engine (which cannot be controlled by users) is used for query processing, and SPARK-based, where DataFrames are created as a virtual data model. To make the results comparable, we choose SPARK-based Squerall and extend it to support Neo4j. We assess the accuracy of OPTIMA in terms of (1) results (accuracy), (2) time, and (3) CPU and memory usage compared to SPARK-based Squerall. We should note that comparing the overall execution time of OPTIMA against an original system, e.g., relational for a given query, is impossible because we are querying various heterogeneous formats and models.
- - In the second study, we inspect OPTIMA's main components: machine learning, data wrappers, and query exe-cution. We observed the behavior of query execution for GRAPH and TABULAR in terms of time. For the data wrapper, we investigate the time taken for the transformation process from data sources to GRAPH or TABU-LAR. As for the machine learning component, we compare our model with the LSTM model in terms of accuracy and time. The LSTM model takes the encoded vectors of SPARQL features as input without any correlation and outputs the data model.

3.2. Experiment 1

In this experiment, we load BSBM* as described above to obtain the results from OPTIMA and SPARK-based Squerall. Then, we run 5150 SPARQL queries and compare the results. This comparison allows us to confirm the correctness of the results returned by OPTIMA. Table 3a shows the results of OPTIMA and SPARK-based Squerall of a complex SPARQL query Q21. The results are the same for both systems, which confirms that OPTIMA is able to support and join large data coming from different datasets.

Table 2 illustrates the execution time returned by both systems. As can be observed, OPTIMA excels Squerall for queries that involve multiple joins. The time difference ranges from 0 to 80000 milliseconds (ms). This difference is due to the predicted virtual data model, e.g., Q19, Q20, in which deep learning predicted that the Virtual model of type GRAPH is optimal. We also observe a small difference in the execution time (ranging from 0 to 30 ms) in favor of Squerall compared to OPTIMA for queries that involve multiple projections, e.g., Q7, Q10. This is explained by the fact that the optimal virtual model is identical to Squerall's, and both Squerall and OPTIMA used the same APIs to call data (wrapper). However, the data model prediction time added to OPTIMA makes it slightly slower than Squerall. Furthermore, the average execution time of Squerall is greater than 4000 ms compared to the average execution time of OPTIMA 2400 ms, as shown in table 3b. These results illustrate the benefits of OPTIMA over existing systems; thus, RQ1 and RQ5 are answered.

N. Bachir Belmehdi et al. / OPTIMA

System OPTIMA	Q1	Q2	Q3	Q4	Q5	Q6	Q7	Q8	Q9	Q10	Q11	Q12	Q13	Q14	Q15	Q16	Q17	Q18	Q19	Q20
Squerall	4098	2519	3091	10299	10199	7984	7089	8427	10094	4684	2561	1400	4644	4469	3885	2875	3314	8742	9059	7407
Time Difference	2807	1265	2361	16	8	6431	15	15	6	10	14	1167	29	18	1488	6	1616	4135	6255	1759
						-			Table	2	• • •	a								
						Time	ın ms	per Qı	iery of	OPTIM	AA &	Squer	all							
To check	if dee	ep lea	arnin	g is re	ducing	g the	overa	all ex	ecutio	n tin	e of	OPT	[MA]	by se	lectii	ng the	e opti	mal v	virtua	l dat
model. We	first i	llustı	rate t	he tim	e take	en by	OPT	IMA	's con	npon	ents:	macl	hine 1	learni	ng al	lgorit	hm, c	query	exec	ution
over GRAP	H mo	odel,	and o	query	execu	tion	over '	ГАВ	ULAR	agai	nst S	PAR	K-ba	sed S	quer	all. V	Ve ru	n OP	TIM	A and
Squerall ov	er 10	30 q	uerie	s. Res	ults a	re sh	nown	in ta	ble 5a	a. Th	e ave	rage	exec	ution	time	e of t	he m	achir	ne lea	rning
component	is ver	y sho	ort, 12	2 ms, •	while	the a	verag	e tim	e for	GRA	PH is	3132	0 ms	and 🛛	ABU	JLAI	R is 2	862 r	ns. R	esult
show that G	RAP	H is	faste	r than	TABU	JLAI	R for	most	queri	es, ev	en w	ith p	redic	tion t	ime.	In su	mma	ry, o	nly 14	1% o
the queries	were	initia	ally fa	aster f	or OP	TIM	A (us	ing (GRAP	H as	a virt	ual n	nodel) con	npare	ed to	Sque	rall a	nd be	cam
in the later	favor.	Thi	s is e	xplain	ed by	the	fact t	hat fo	or thos	se qu	eries,	, ther	e is a	ı slig	nt dif	feren	ice in	exec	utior	time
using GRA	PH cc	mpa	red to	o Sque	erall. T	This a	answe	ers R	Q2.											
Finally, v	ve rec	cord	the F	Resour	ce Co	onsun	nptio	n (i.e	., Me	mory	and	CPU) tak	en by	OP'	TIM	A and	I SPA	RK-	based
Squerall. T	he res	sults	repo	rted in	ı Tabl	e 3c	show	/ that	t the C	CPU	is no	t full	ly use	ed by	OP'	TIM	A and	SPA	RK-	based
Squerall (ar	ound	0.21	% w	as use	d). Th	is m	eans	that t	he coi	mple	kity o	of que	eries	does	not i	mpac	t CP	U coi	nsum	ption
As for the to	otal m	emo	ry res	served	, OPT	IMA	cons	umec	l abou	t IGI	3 ove	r 8G	B per	node	, whi	ile SF	ARK	-base	ed Sq	ueral
used at mos	t 1GI	3. Th	ie sai	ne CF	U and	1 mei	mory	coul	d be e	xplai	ned b	by the	e fact	that	both	are u	ising	the s	ame	quer
engine - SP.	ARK	and t	the di	stribu	tion o	f CPI	U bet	ween	the ne	odes	for lo	adin	g and	trans	sform	natior	ı. Thi	s ans	wers	RQ7
e													0							-
Query SELECT DISTINCT	20roduotI aba	1 2producar	OP	ΓIMA		S	Sqyerall			Sy	stem	Т	ïme (n	ns)						
WHERE { product re 2producer rdfs:label	fs:label ?prod	uctLabel .	['B	r Mix Lemon'	Coke Classic 3	55 MP1 F	'Bar Mix Len	ion' 'Coke C	lassic 355 MI'l	OI	TIMA	1 24	400							
?product rdf:type bsb ?product bsbm:produ	m:Product . cer ?producer	.}	1.1.1						,	Sq	uerall	4	200							
(a) table	3010	llory	Decul	t Datu	mad by		гіма	8 80	uaroll				<i>a</i> >	. 11	21 4	т.				
(a) table	Ja. Q	uciy	Resul		neu og			æ sy	ucran				(0)	table	30: A	Ng II	me			
					N	Aetric	s	0	PTIM	Α	Sque	erall								
					CPU	averag	ge (%)	0.	21		0.20									
					Max n	nemor	y (GB)	1	.0		0.97									
						(c)) table	Sc. F	Resour	ce Co	nsumi	ntion								
						(0)) their		cobo un			puon								
2 2 F		•																		
3.3. Experi	ment	2																		
· ··								-	0 D											
In this stu	idy, w	ve ev	aluat	e the r	nain c	comp	onent	s of (OPTIN	MA.										
3.3.1. Maci	hine I	Learn	ing (Compe	onent															
We evalu	ated of	our n	nodel	with	an LS	TM	mode	l to a	issess	our e	encod	ling t	echni	iaues	and	predi	ction	mod	lel. W	e us
5150 auerie	s. 80 ^o	% are	euseo	d for t	rainin	g, and	d 20%	6 are	used f	for v	lidat	ion. V	We tr	ain b	oth m	nodel	s on t	he sa	me d	atase
and comput	e the	accu	racv	and C	ross-e	ntron	v los	s fun	ction	Next	. we	evalu	ate th	ne mo	dels	' effic	iency	in te	erms	of the
models' tra	ining	time	and	predic	tion t	ime.	Table	e 5b	shows	s that	our	tree-s	struct	ured	base	d me	hod o	outpe	rforn	ns the
LSTM mod	el wit	h an	avera	ige acc	curacy	of 0	.831 f	or or	ir mod	lel ag	ainst	0.70	8 for 1	the L	STM	mod	el. Th	e cro	ss-en	trop
loss is equa	l to (0.000	18 fc	or our	mode	l cor	nnare	d to	1 920	27 fc	r LS	тм	This	is be	caus	e the	LST	M m	odel	relie
			~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~ ~	~ ~ ~ 1		лсол	IIDan		1.720	21		1 1 1 1 1	110.5	10	~~~				UQU-	
on the inde	bende	nt as	sum	otion a	mono	diffe	erent	oper	ations	and	attrib	utes	while	e tree	-stru	cture	d-bas	ed m	ethor	ls cai
on the indep capture the	oende corre	nt as latio	sump ns be	otion a tween	mong	diffe	erent s and	opera attril	ations	and a	attrib es. O	utes, ur m	while odel :	e tree	-stru ves f	cture he be	d-bas	ed m	ethoc	ls cai

3.3.2. Query Execution As shown in Table 4, the analysis of experimental results indicates that GRAPH is faster than TABULAR in most cases, except for queries like Q8 and Q10. It has comparable to slightly lower performance in Q16. This

	Q1	Q2	Q3	Q4	Q5	Q6	Q7	Q8	Q9	Q10	Q11	Q12	Q13	Q14	Q15	Q16	Q17	Q18	Q19	Q20
Prediction Time	3	3	4	6	4	5	5	6	2	4	5	1	5	5	4	3	2	4	4	4
GRAPH	1143	1161	1239	1243	306	3181	7168	12237	4977	16681	1211	3567	482	1285	766	2883	6639	1366	3370	1723
TABULAR	4098	2519	3091	10283	10191	7984	7089	8427	10088	4684	2561	1400	4644	4469	3885	2875	3314	8742	9059	7407
					Time in	n ms j	per Qı	uery of	Predict	tion, Gl	RAPH	& TA	BULA	R						
	(Condit	ion	A	vg. time	(ms)				Cos	t	Lo	88	Асси	racy]				
	Mach	nine Le	earnin	g	12					LOT	- M	1.02	227	0.7	00					
	On	ly GR	APH		1320		1		_	LSII	.vi	1.92	JZ7	0.7	08					
	Only	y TAB	ULAR	2	2862					Our Mo	odel	0.00	018	0.8	31					
									(1	\ m 11		-			c	D	т			

confirms that the optimal model is very important to reduce the execution time of queries. The total execution time ranges from 50 to 90000 ms, with 90% of all cases being about or below 3000 ms. OPTIMA virtual data model of type GRAPH is faster in queries that involve joins (ranging from 50 to 40000 ms), while the TABULAR model outperforms the GRAPH model in queries involving more projections (ranging from 200 to 90000 ms).

This is explained by the fact that the GRAPH is designed to store connections between data. Therefore, queries do not scan the entire graph to find the nodes that meet the search criteria. It looks only at nodes that are directly connected to other nodes, while SQL-like methods used by the TABULAR model require expensive join operations because they traverse all data to find the data that meets the search criteria. On the other hand, the TABULAR model is faster when handling projections because the data structure is already known, and data can be easily accessed by column names. Conversely, the GRAPH model does not have a predefined structure for the data, and each node attribute has to be examined individually during the projection query.

The number of joins has a decisive impact on query performance; it should be taken into consideration with other factors, e.g., size of involved data, presence of filters, and selected variables. For example, Q2 joins only two data sources, Product and Review (1254 ms) but has comparable performance with Q1 (1291 ms), which joins four entities (Product, Offer, Review, and Producer). This may be due to the presence of filtering in Q1 (16 filters), which significantly reduces intermediate results to join. Q3 involves four data sources, yet it is among the fastest queries. This is because it involves the small entities Person and Producer, which is another reason to reduce intermediate results to join. With five data sources to join, Q4 is among the most expensive queries (10299 ms). This can be attributed to the fact that the filter on Product is selective (?language = "en"), which results in large intermediate results to join, in contrast to Q6 (? price < 8000). Although the four-source join Q7 and Q8 involve the small entity Producer, they are the most expensive queries that execute over the GRAPH model; this can be attributed to many projections (38 attributes). Thus, we answer RQ3 and suggest that operations can affect query execution time.

Model	Neo4j	JDBC	CSV	Cassandra	MongoDB	Loading
GRAPH	138	954	196	7695	188	4.327
TABULAR	3275	199	255	5319	330	7.141
			Tabl	e 6		

Time (ms) of Data transformation to GRAPH & TABULAR

3.3.3. Data wrapper Time

To answer RQ8, we evaluate, in this study, the time needed to load the data from data sources to the virtual data model of type GRAPH or TABULAR (see Table 6). Since the transformation process is different, we expect different wrapper behavior. In the table, we illustrate the time needed by each wrapper with the following observations:

 Neo4j connector loads 50000 nodes from Neo4j within 138 ms into GRAPH, compared to 3275 ms in TABULAR. This is explained by the fact that the graph property used by Neo4j has the same structure as the GRAPH model.

- CSV connector loads 50000 rows within 196 ms from CSV files into GRAPH, compared to 255 ms in TABULAR, even though CSV files save data into tables. This can be explained by the fact that the GRAPH virtual model is a schema-less model that loads data directly without the need to preserve data structure, while TABULAR takes time to build the data schema.
- JDBC connector loads 50000 rows from MySQL database within 954 ms into GRAPH, compared to 199 ms in TABULAR. This can be explained by the fact that MySQL uses a relational model with the same data structure as the Virtual TABULAR model.
- MongoDB connector loads 50000 rows from MongoDB within 188 ms into GRAPH, compared to 330 ms in TABULAR. This can be explained by the fact that MongoDB is document-based, i.e., schema-less, the same as the GRAPH Virtual model, unlike the TABULAR model, which needs to build a data schema.
- Cassandra connector loads 50000 rows within 7695 ms into GRAPH, compared to 5319 ms in TABULAR. This can be explained by the fact that Cassandra uses a columnar data model, which is closer to the TABULAR model, even though it is a NoSQL database.

4. Related Work

Our literature review reveals two categories addressing data virtualization. These two categories are namely "ontology-based data access" and "non-ontology-based data access" [6]. Non-ontology-based data access approaches mostly use SQL-like as query language and implement a virtual relational model [14, 15], defining views of relevant data from sources having a relational model. Those views are generated based on mapping assertions that associate the general relational schema with the data source schemata. The shortcomings of these approaches is that the schema modifications and extensions are very rigid due to mappings and may depend on complex constraints. Furthermore, these approaches use Self-Contained Query [16] where users cannot control the structure of the virtual data model. OBDA [17] approaches use SPARQL as a unified access language and detect relevant data from sources to be joined through ontologies and standardized mappings. This provides flexibility in modifying and extending the ontology and mappings with semantic differences found across the data schemata.

We identified commercial systems such as Stardog (www.stardog.com), and Oracle Spatial and Graph (www. oracle.com/technetwork/database/options/spatialandgraph, Mastro [18], Ultrawrap [19] and open-source systems such Ontop [3], Morph [20], which implemented OBDA by querying relational data sources using virtual knowledge graph. These solutions, however, are not designed to query large-scale data sources, e.g., NoSQL stores or HDFS.

Our study's scope focuses on works that query large-scale data sources using OBDA. Optique [4] is an OBDA platform that accesses both static and streaming data. It implements a relational model (implicitly a TABULAR) as a virtual model while querying data sources such as SQL databases and other sources, e.g., CSV, and XML. There was no clear description of how Optique accesses NoSQL stores and distributed file systems (e.g., HDFS). Ontario [5] focuses on query rewriting, planning, and federation, with a strong stress on RDF data as input. Query plans are built and optimized based on a set of heuristics. The virtual model used by Ontario is the GRAPH model (explicitly an RDF). Squerall [6], a recent and close work to OPTIMA, leverages Big Data engines SAPRK and Presto to query on-the-fly of heterogeneous large data sources. The virtual data model imposed by Presto is TABULAR and does not offer users to control it; while SPARK can offer control over the virtual data model, Squerall uses DataFrame as a virtual model, which is TABULAR. However, the decision behind the virtual data model implemented by all these systems is based on use and flexibility and not on solid evidence to improve query processing. No work exists that (1) implements the different optimal virtual models and (2) selects the optimal one based on query behavior. For machine learning, some works [21–23] addressed the cost estimation of SPARQL queries to optimize query execution plan, e.g., performance prediction, however, all these approaches are designed for a single query on one single data source.

5. Conclusion

We implemented OPTIMA - an ontology-based big data access system that reduces query time execution by predicting the optimal virtual data model, GRAPH or TABULAR, based on query behavior. The effective deep learning

model built on top of OPTIMA's architecture extracts significant features such as the query plan and operations and predicts the optimal virtual data model that has the lowest query execution time. The experiment showed a reduction in query execution time of over 40% for the TABULAR model and over 30% for the GRAPH model selection. References [1] N. Miloslavskaya and A. Tolstoy, Big data, fast data and data lake concepts, Procedia Computer Science 88 (2016), 300–305. [2] A. Poggi, D. Lembo, D. Calvanese, G. De Giacomo, M. Lenzerini and R. Rosati, Linking data to ontologies, in: Journal on Data Semantics X, Springer, 2008. [3] D. Calvanese, B. Cogrel, S. Komla-Ebri, R. Kontchakov, D. Lanti, M. Rezk, M. Rodriguez-Muro and G. Xiao, Ontop: Answering SPARQL queries over relational databases, Semantic Web 8(3) (2017), 471-487. [4] M. Giese, A. Soylu, G. Vega-Gorgojo, A. Waaler, P. Haase, E. Jiménez-Ruiz, D. Lanti, M. Rezk, G. Xiao, Ö. Özçep et al., Optique: Zooming in on big data. Computer 48(3) (2015), 60-67. [5] K.M. Endris, P.D. Rohde, M.-E. Vidal and S. Auer, Ontario: Federated query processing against a semantic data lake, in: International Conference on Database and Expert Systems Applications, Springer, 2019, pp. 379–395. [6] M.N. Mami, D. Graux, S. Scerri, H. Jabeen, S. Auer and J. Lehman, Squerall: Virtual Ontology-Based Access to Heterogeneous and Large Data Sources, Proceedings of 18th International Semantic Web Conference (2019). [7] J.M. Hellerstein, Optimization techniques for queries with expensive methods, ACM Transactions on Database Systems (TODS) 23(2) (1998), 113-157. [8] S. Batra and C. Tyagi, Comparative analysis of relational and graph databases, International Journal of Soft Computing and Engineering (IJSCE) 2(2) (2012), 509-512. [9] C.B.B. Belmehdi, A. Khiat and N. Keskes, OPTIMA: Framework Selecting Optimal Virtual Model to Query Large Heterogeneous Data, in: Big Data Analytics and Knowledge Discovery - 24th International Conference, DaWaK 2022, Vienna, Austria, August 22-24, 2022, Proceedings, Vol. 13428, 2022, pp. 209-215. [10] M. Hunger, From Relational to Graph: A Developer's Guide, DZone, ed (2016). [11] M.N. Mami, S. Scerri, S. Auer and M.-E. Vidal, Towards Semantification of Big Data Technology, in: Big Data Analytics and Knowledge Discovery, S. Madria and T. Hara, eds, Springer International Publishing, Cham, 2016, pp. 376-390. ISBN 978-3-319-43946-4. [12] M. Atay, Y. Sun, D. Liu, S. Lu and F. Fotouhi, Mapping XML data to relational data: A DOM-based approach, arXiv preprint arXiv:1010.1746 (2010). [13] J.E. Gonzalez, R.S. Xin, A. Dave, D. Crankshaw, M.J. Franklin and I. Stoica, {GraphX}: Graph Processing in a Distributed Dataflow Framework, in: 11th USENIX symposium on operating systems design and implementation (OSDI 14), 2014, pp. 599-613. [14] R.F. van der Lans, Architecting the multi-purpose data lake with data virtualization, Denodo whitepapers (2018). [15] D. Chatziantoniou and V. Kantere, Just-In-Time Modeling with DataMingler, in: Proceedings of the ER Demos and Posters 2021 co-located with 40th International Conference on Conceptual Modeling (ER 2021), St. John's, NL, Canada, October 18-21, 2021, R. Lukyanenko, B.M. Samuel and A. Sturm, eds, CEUR Workshop Proceedings, Vol. 2958, CEUR-WS.org, 2021, pp. 43-48. http://ceur-ws.org/Vol-2958/ paper8.pdf. [16] M.N. Mami, D. Graux, S. Scerri, H. Jabeen, S. Auer and J. Lehmann, Uniform Access to Multiform Data Lakes using Semantic Technolo-gies, in: Proceedings of the 21st International Conference on Information Integration and Web-based Applications & Services, iiWAS 2019, Munich, Germany, December 2-4, 2019, ACM, 2019, pp. 313–322. doi:10.1145/3366030.3366054. [17] D. Calvanese, G.D. Giacomo, D. Lembo, M. Lenzerini, A. Poggi, M. Rodriguez-Muro and R. Rosati, Ontologies and databases: The DL-Lite approach, in: Reasoning Web International Summer School, Springer, 2009, pp. 255-356. [18] D. Calvanese, G. De Giacomo, D. Lembo, M. Lenzerini, A. Poggi, M. Rodriguez-Muro, R. Rosati, M. Ruzzi and D.F. Savo, The MASTRO system for ontology-based data access, Semantic Web 2(1) (2011), 43-53. [19] J.F. Sequeda and D.P. Miranker, Ultrawrap: SPARQL execution on relational data, Journal of Web Semantics 22 (2013), 19–39. [20] F. Priyatna, O. Corcho and J. Sequeda, Formalisation and Experiences of R2RML-Based SPARQL to SQL Query Translation Using Morph, in: Proceedings of the 23rd International Conference on World Wide Web, 2014, pp. 479-490-. [21] W.E. Zhang, Q.Z. Sheng, Y. Qin, K. Taylor and L. Yao, Learning-based SPARQL query performance modeling and prediction, World Wide Web 21(4) (2018), 1015–1035. https://doi.org/10.1007/s11280-017-0498-1. [22] R. Hasan and F. Gandon, A machine learning approach to sparql query performance prediction, in: 2014 IEEE/WIC/ACM International Joint Conferences on Web Intelligence (WI) and Intelligent Agent Technologies (IAT), Vol. 1, IEEE, 2014, pp. 266-273. [23] R. Singh, Inductive learning-based sparql query optimization, in: Data Science and Intelligent Applications, Springer, 2021, pp. 121–135.